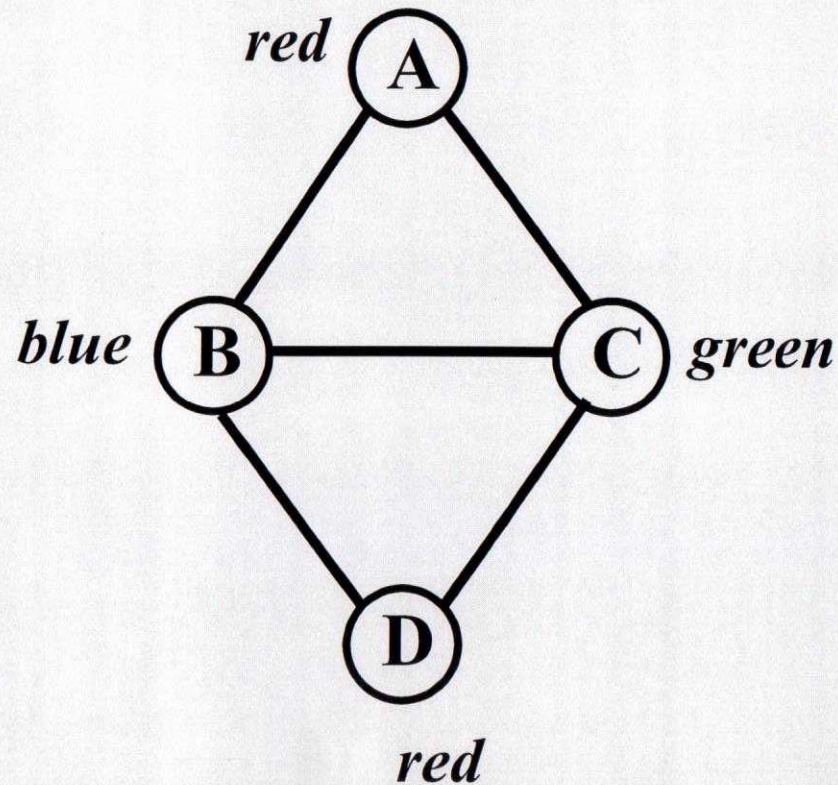


Constraint satisfaction problems (CSPs)

- Defined by:
 - A set of variables x_1, x_2, \dots, x_n
 - A domain D_i for each variable x_i
 - Constraints c_1, c_2, \dots, c_m
- A constraint is specified by
 - A subset (often, two) of the variables
 - All the allowable joint assignments to those variables
- Goal: find a complete, consistent assignment

Graph coloring

- Fixed number of colors; no two adjacent nodes can share a color



Cryptarithmic puzzles

$$\begin{array}{r} \text{TWO} \\ \text{TWO} + \\ \hline \text{FOUR} \end{array}$$

E.g., setting $F = 1$, $O = 4$, $R = 8$, $T = 7$, $W = 3$,
 $U = 6$ gives $734 + 734 = 1468$

Cryptarithmic puzzles...

T W O

Trick: introduce auxiliary

T W O +

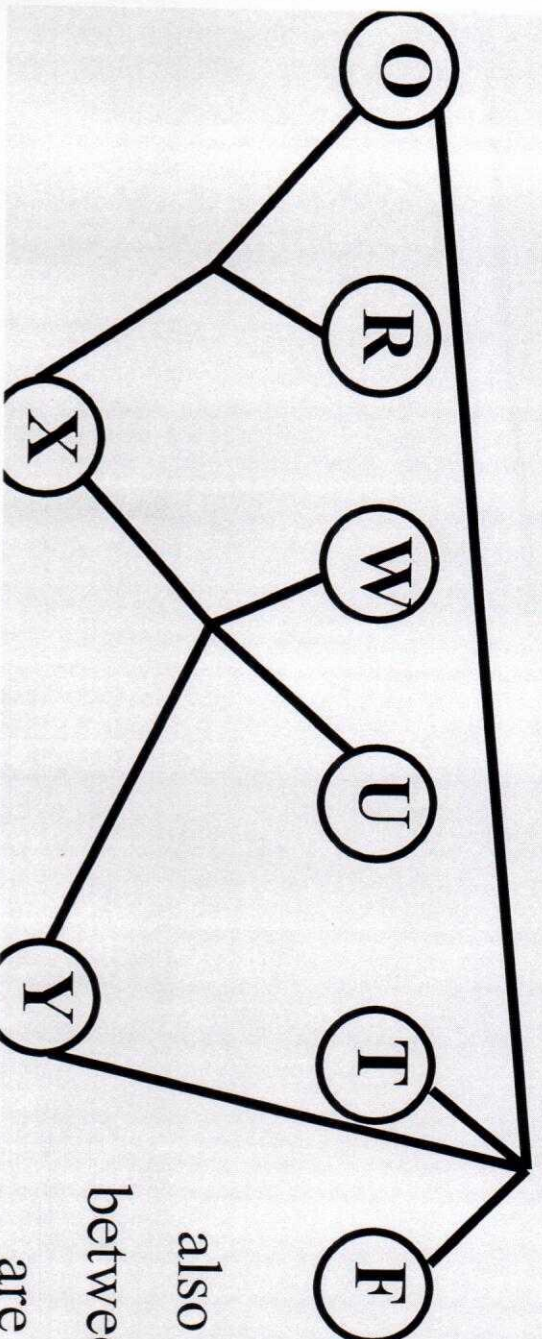
variables X, Y

FOUR

$$O + O = 10X + R$$

$$W + W + X = 10Y + U$$

$$T + T + Y = 10F + O$$



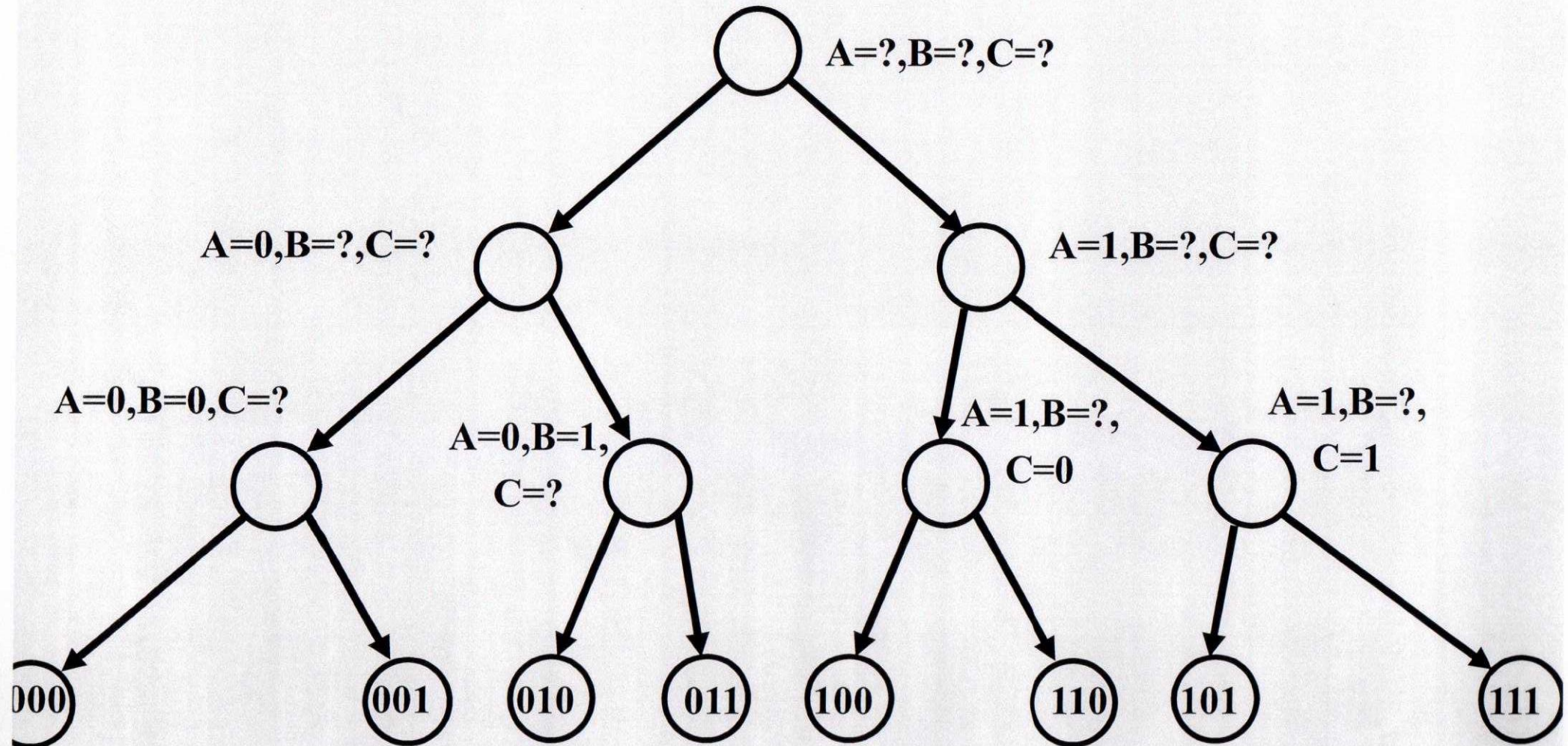
also need pairwise constraints
between original variables if they
are supposed to be different

Generic approaches to solving CSPs

- State: some variables assigned, others not assigned
- Naïve successors definition: any way of assigning a value to an unassigned variable results in a successor
 - Can check for consistency when expanding
 - How many leaves do we get in the worst case?
- CSPs satisfy commutativity: order in which actions applied does not matter
- Better idea: only consider assignments for a single variable at a time
 - How many leaves?

Choice of variable to branch on is still flexible!

- Do not always need to choose same variable at same level
- Each of variables A, B, C takes values in $\{0,1\}$



- Can you prove that this never increases the size of the tree?